

CSCI 210: Computer Architecture

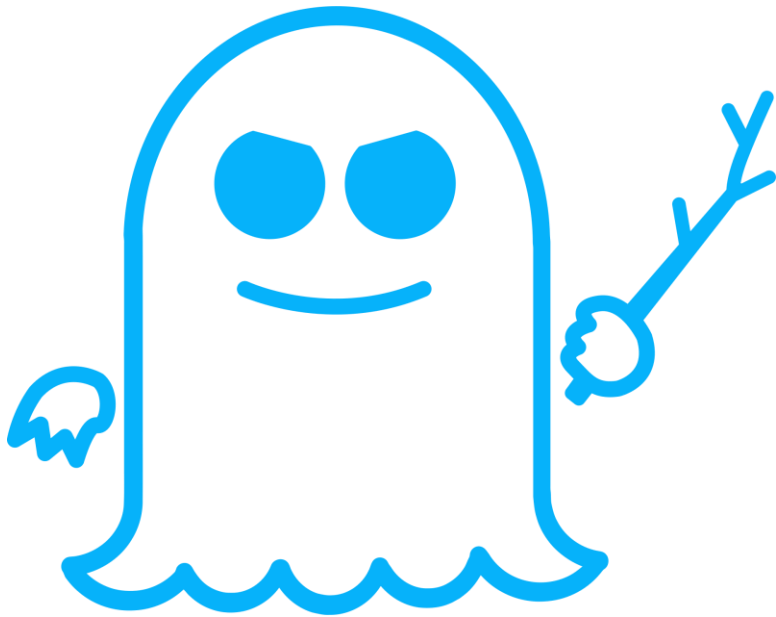
Lecture 37: Last Class!

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CS History: Cache Side Channel Attacks

- Flush + Reload (2014), Evict + Reload (2015)
- Depends on shared memory (e.g., a shared crypto library) being in the cache
- Attacker causes the cache to be emptied, either by calling flush or by filling it with new data
- Victim uses shared memory to do something (e.g., encrypt a message)
- Attacker references different bytes of shared memory and measures how long it takes: if it's fast, byte is in the cache, which means the victim referenced it. If it's slow, not in the cache, so victim didn't reference it
- This lets the attacker trace the code the victim is running
- Researchers used it to extract private encryption keys, even across virtual machines
- Takes advantage of L3 cache being shared across cores!

CS History: Spectre & Meltdown (2018)



SPECTRE

- Combines speculative execution with cache side channel attacks
- Create a program with a pattern of references that will cause the branch predictor to predict taken
- Put some code you don't have access to in the branch – it will be run by the branch predictor, but then rolled back
- Now the memory it referenced is in the cache!
- Use cache side channel attack to trace the code
- Meltdown does a similar thing, but accessing memory in the OS kernel
- Affects every Intel processor since 1995, also ARM processors, all OSes
- Fixes aren't cheap and tend to slow processors down a lot

Cache Final Project

- Double check that the skeleton code provided includes support for the command line flag `--hit-time` (or `-h`) for setting the cache hit time
 - If it doesn't apply the patch in my BlackBoard announcement

Options:

<code>-b, --block-size <BLOCK_SIZE></code>	Set the block size in bytes. Must be a power of two [default: 32]
<code>-a, --associativity <ASSOCIATIVITY></code>	Set the associativity. Must be a power of two [default: 1]
<code>-s, --size <SIZE></code>	Set the cache size in kilobytes (kB). Must be a power of two [default: 32]
<code>-h, --hit-time <HIT_TIME></code>	Set the cache hit time [default: 4]
<code>-m, --miss-penalty <MISS_PENALTY></code>	Set the cache miss penalty [default: 34]
<code>--help</code>	Print help
<code>-V, --version</code>	Print version

Questions on Cache Final Project?

We have a 8 byte block size, direct-mapped, 16 kB cache. For a given byte address, to find the offset

- A. Mod by 8
- B. Mod by 16
- C. Divide by 8
- D. Divide by 16
- E. None of the above

If we are simulating a cache and not actually accessing data, do we need the offset?

- A. Yes, we need it to access the right cache block
- B. No, we can ignore it entirely and use those bits as part of the index
- C. No, but we still need to divide by the block size to get the block address

We have an 8 byte block size, direct-mapped, 16 kB cache. How many rows will this cache have?

A. 16

*Recall a kB is 1024 bytes

B. $2 * 1024$

C. $8 * 1024$

D. $16 * 1024$

We have an 8 byte block size, direct-mapped, 16 kB cache. How can we find the index from an address?

- A. Divide by 8
- B. Divide by $(16 * 1024)$
- C. Mod by $16 * 1024$
- D. Divide by 8, then mod by $(2 * 1024)$
- E. None of the above

We have an 8 byte block size, 2-way set associative, 16 kB cache. How can we find the index from an address?

- A. Divide by 8, then mod by 1024
- B. Divide by 8, then mod by $(2 * 1024)$
- C. Divide by 8, then mod by $(2 * 2 * 1024)$
- D. Divide by 16, then mod by $(2 * 1024)$
- E. None of the above

We have an 8 byte block size, direct mapped, 16 kB cache. How can we find the tag of an address?

A. Divide by 8

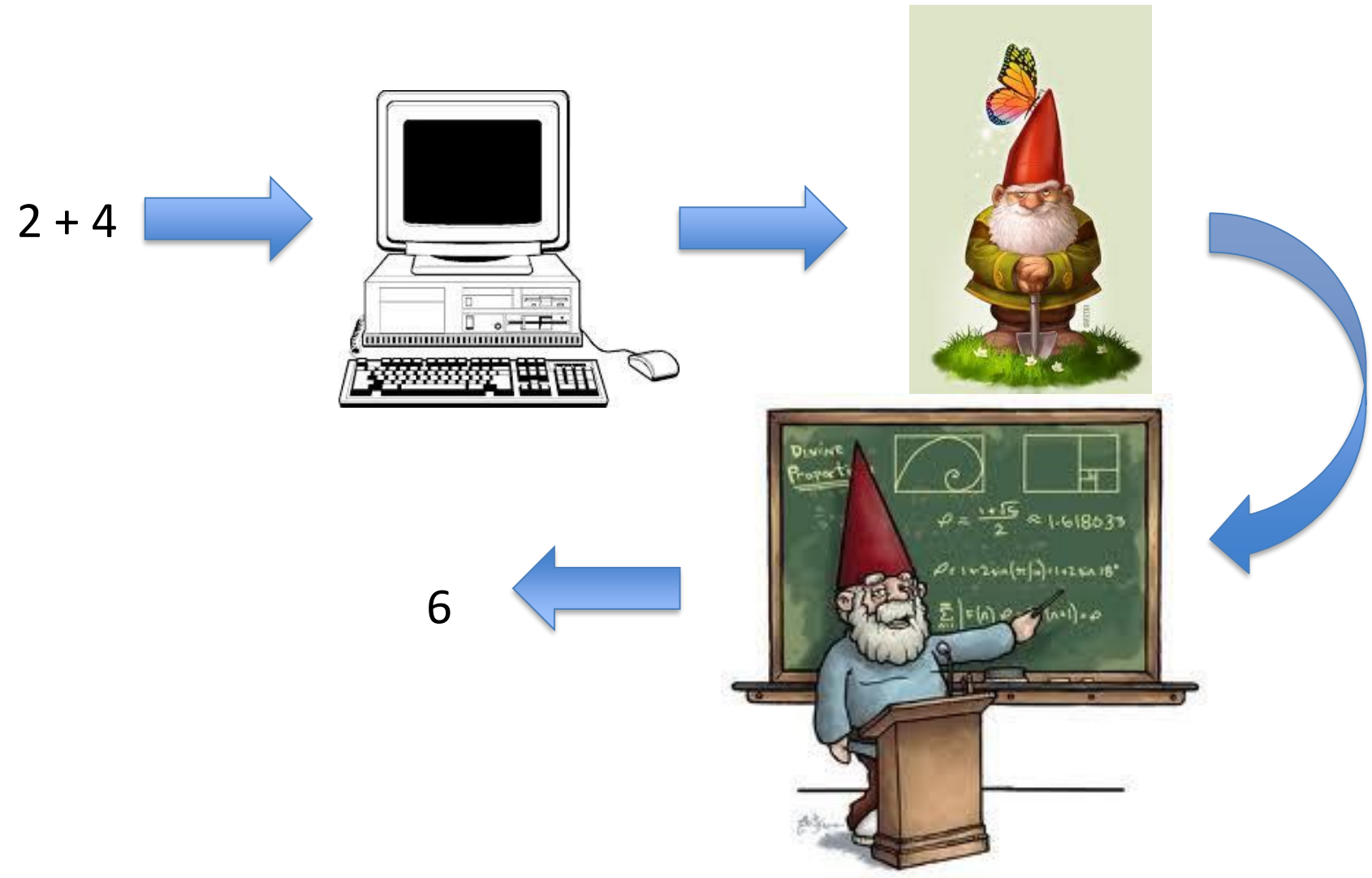
B. Divide by $8 * 1024$

C. Divide by $16 * 1024$

D. None of the above

Questions on Final Project?

Previous Conceptions of How Computers Work



Actually Assembly

High Level:

$x = 2 + 4$

Assembly (assuming we have a mem address for x in \$s0):

```
li      $t1, 2
addi   $t1, $t1, 4
sw     $t1, 0($s0)
```

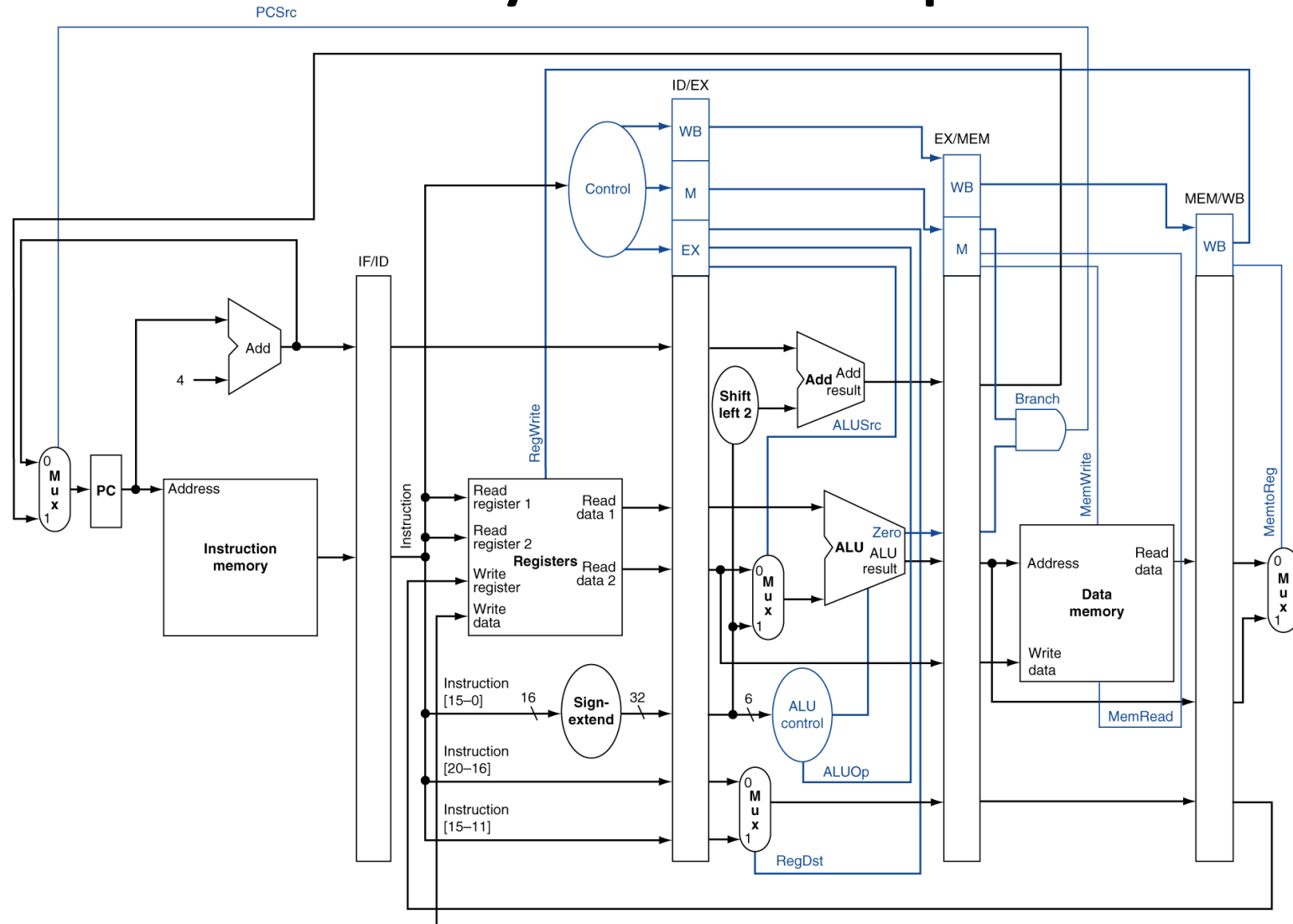
Actually Machine Instructions

```
addi $t1, $t1 5
```

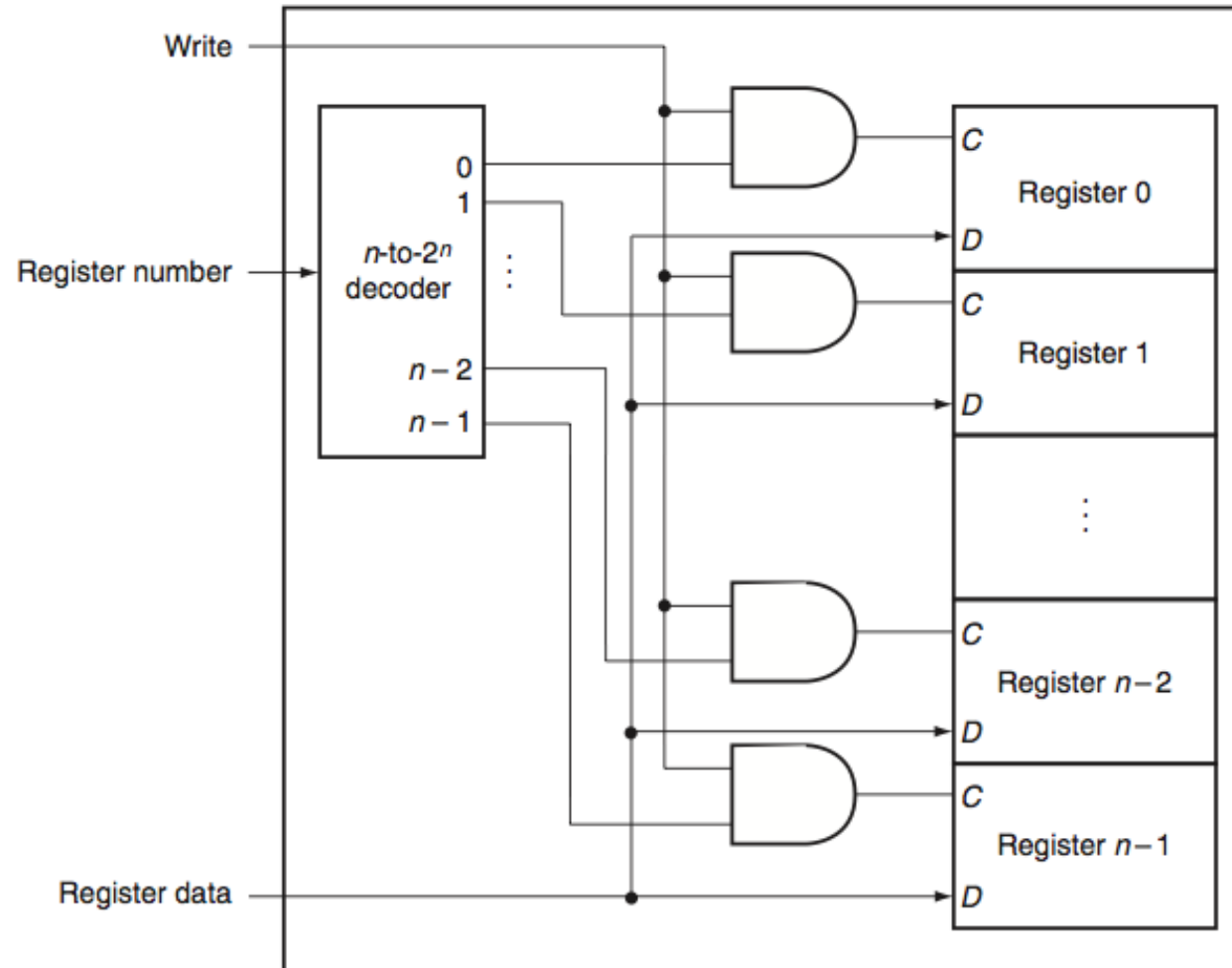


```
0010000100101001000000000000001001
```

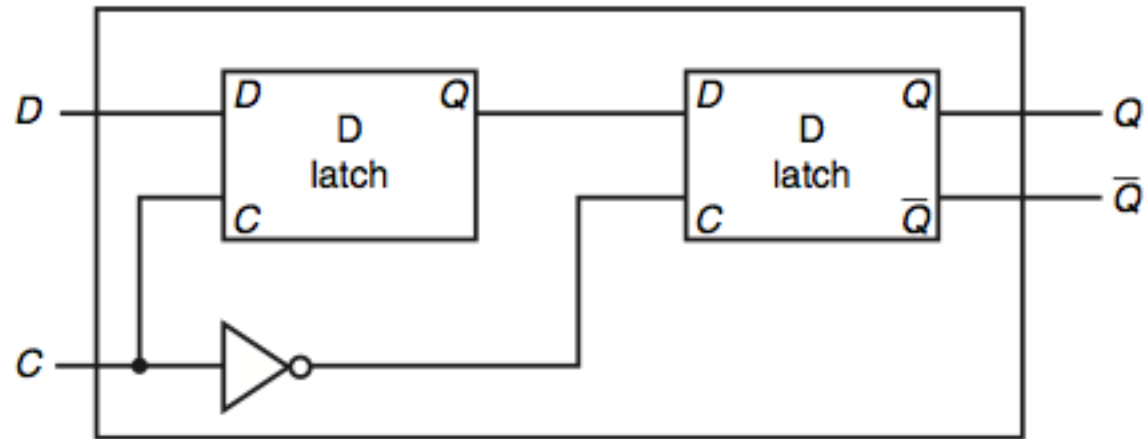
Actually The Datapath



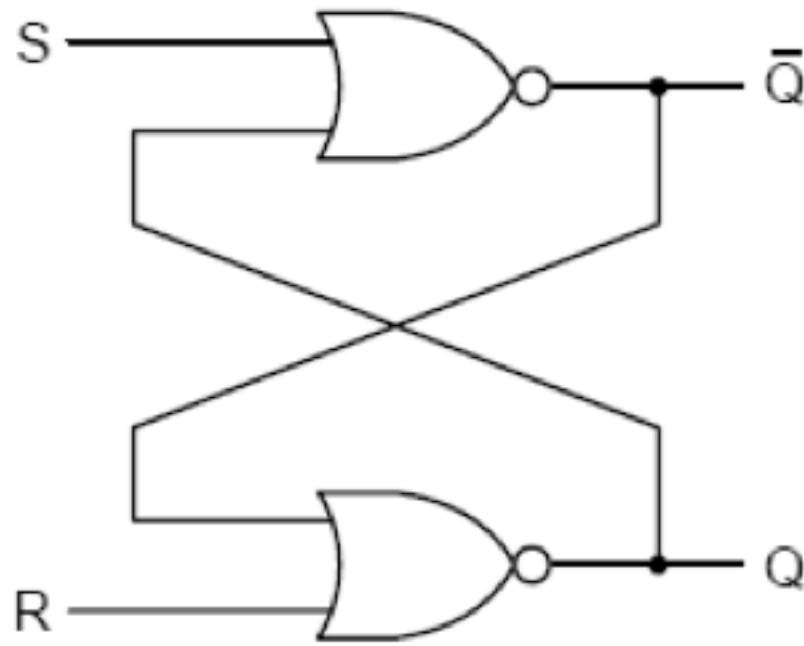
Actually Registers



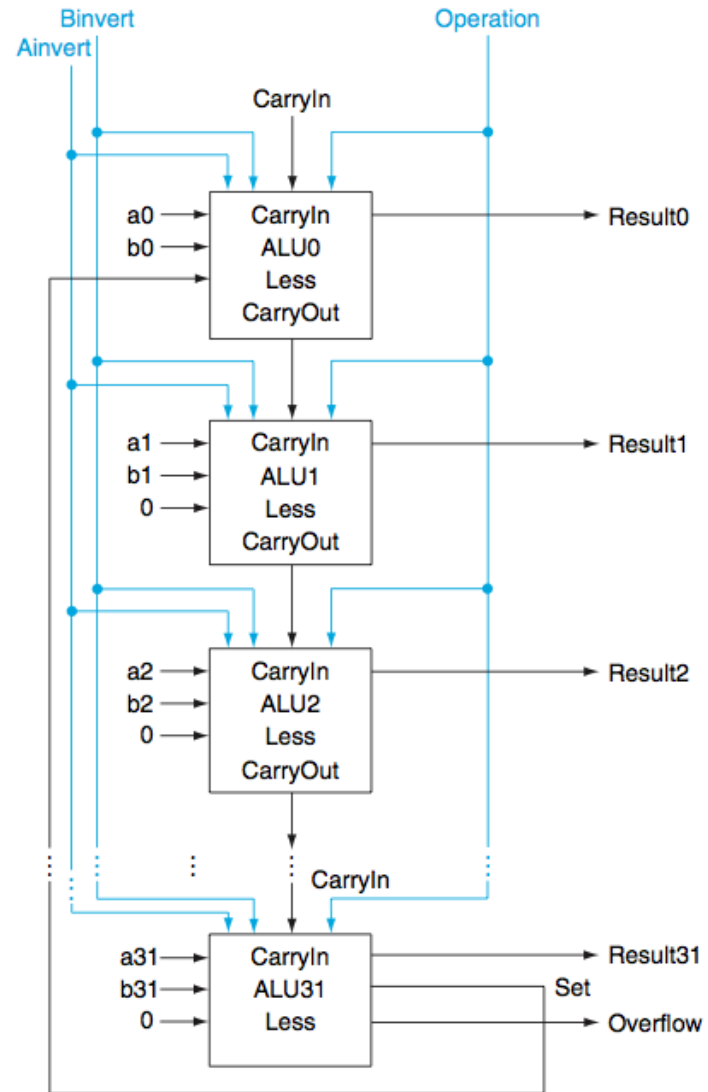
Actually Flip-flops



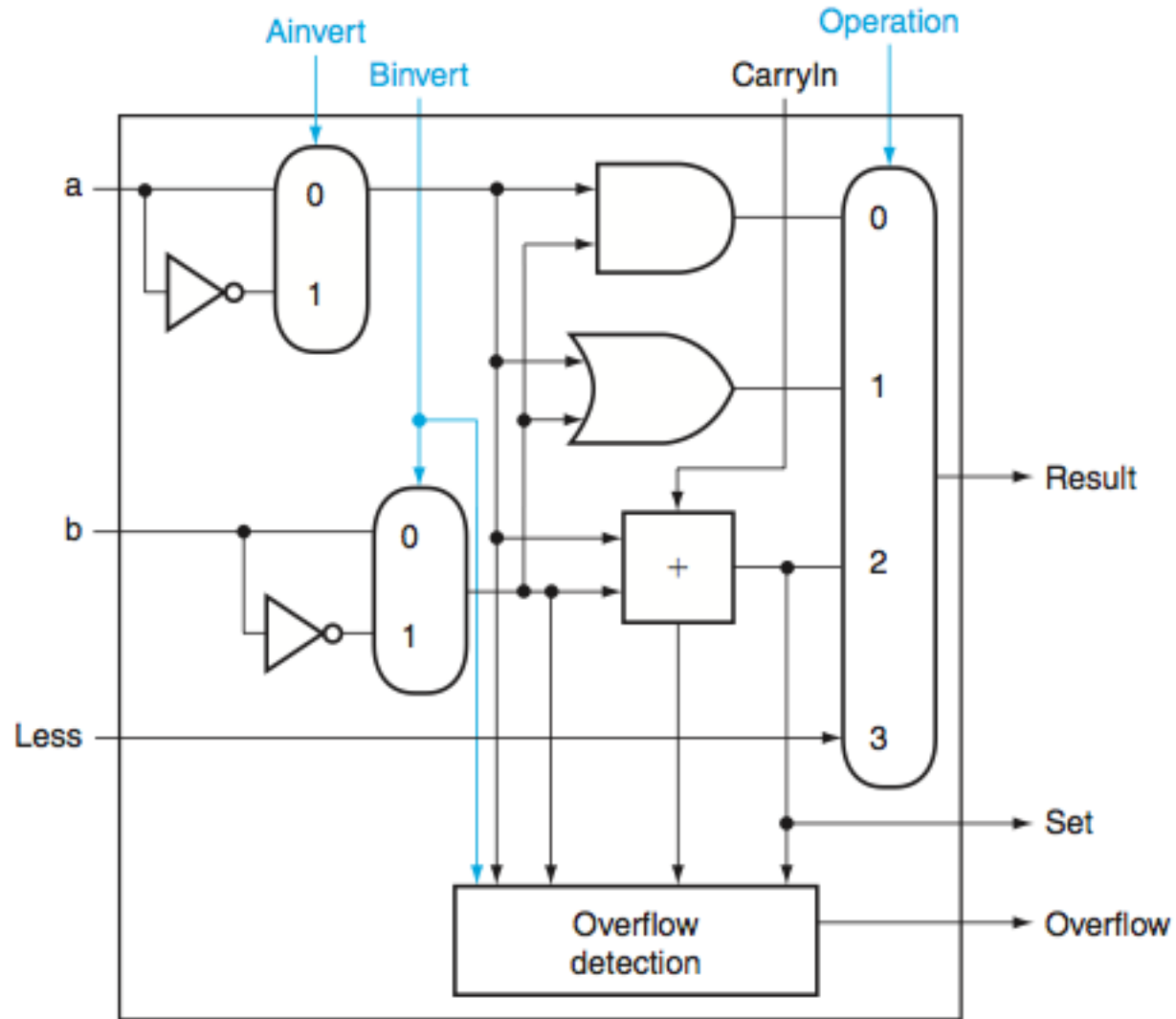
Actually Latches



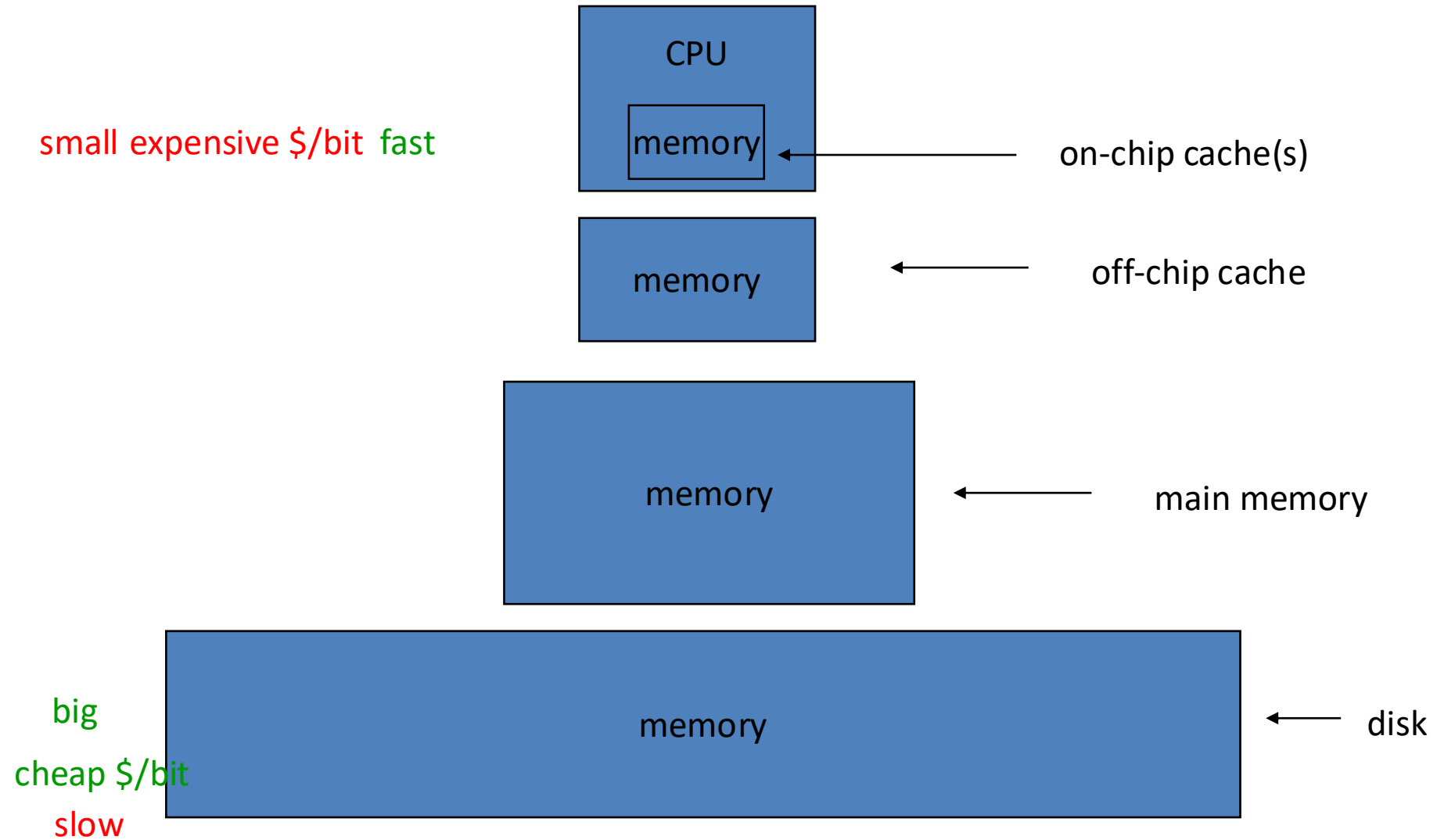
Actually the ALU



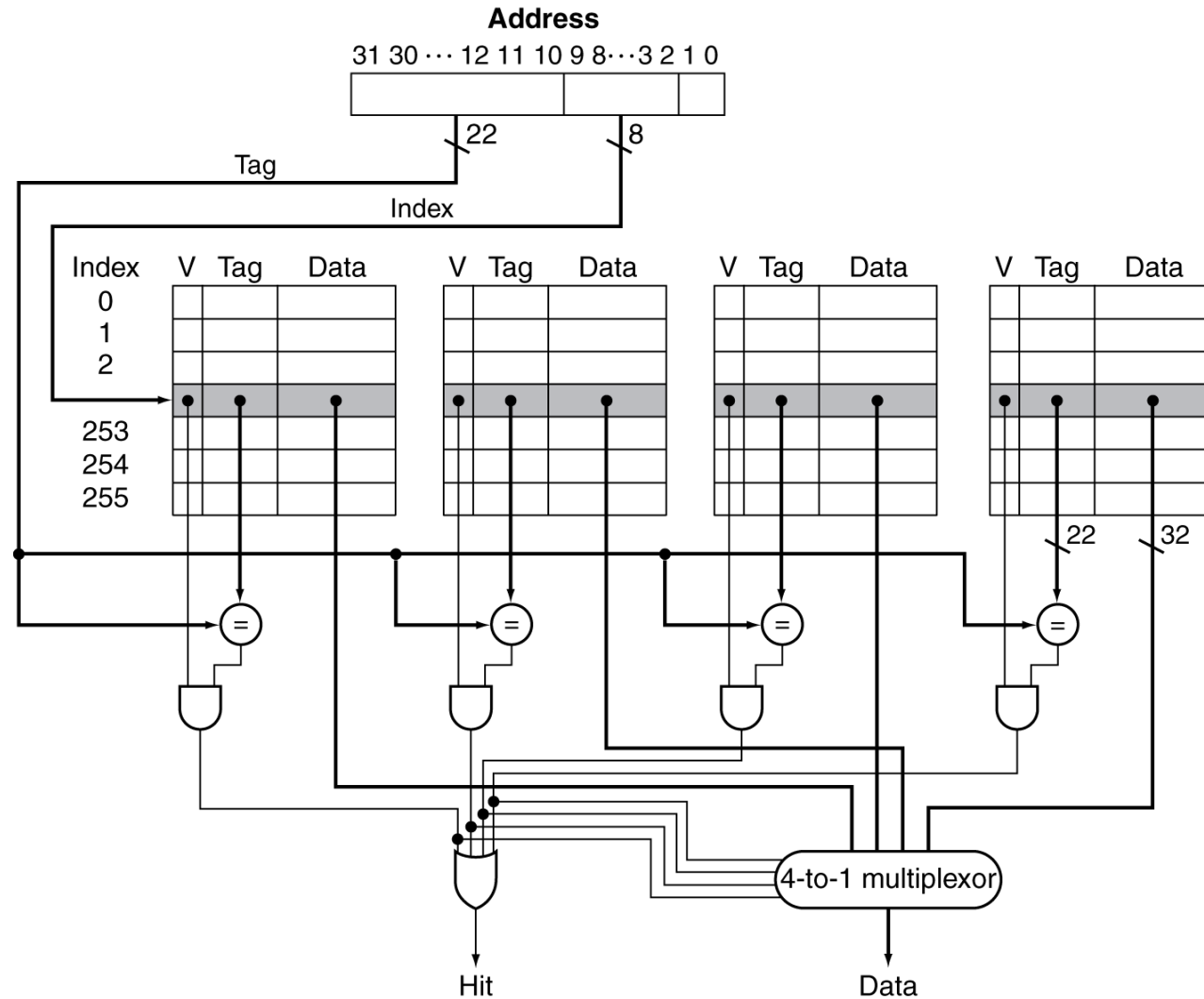
Actually the ALU



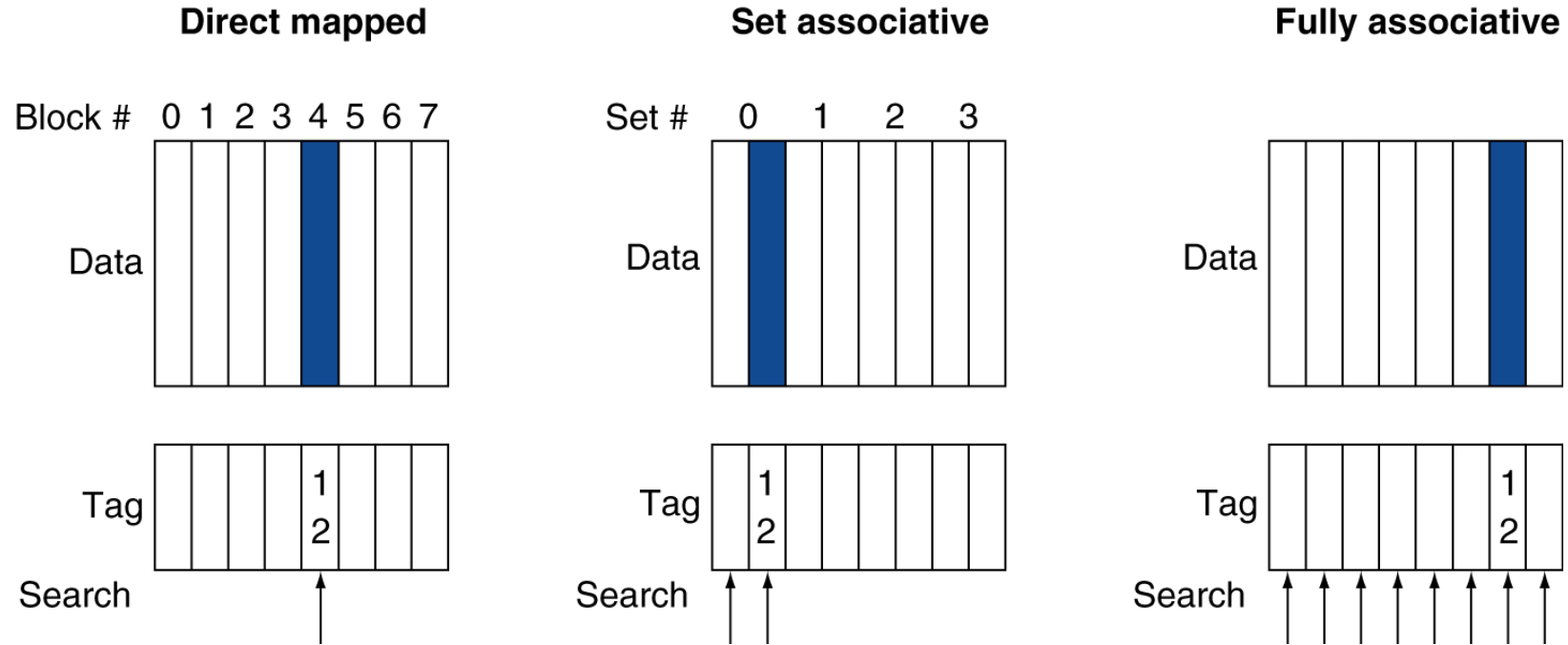
Actually Memory



Actually Caches



Actually LOTS of Caches



But wait, what about?

- Negative Numbers
- Floating Point
- All that other stuff . . .

Computers

are

Complicated

- But now, you know how they work. Kinda.

- I appreciate all the hard work you've done for this class.
- Have a great break!
- ...and fill out course evals!